Writing Fast Haskell

Elegance Is Not an Excuse for Bad Performance

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August 22, 2018

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- But elegance is *not* an excuse for bad performance!
- Writing fast Haskell requires some understanding of GHC's internals
- GHC provides a surprising number of tools to influence performance

Goals for Today

- 1. Learn to reason about performance
- 2. Look under the hood of GHC (specifically Core)
- 3. Learn about some rules of thumb for writing fast Haskell
- 4. Learn about primitives and libraries useful for writing fast Haskell

Benchmarking/Profiling Disclaimer

- Benchmark before you optimize
- GHC supports options for time and space profiling
- Profiling can break optimizations
 - Enable profiling selectively

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Represented by a pointer to a heap object

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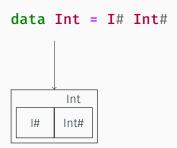
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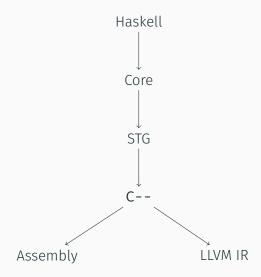
Unlifted Types

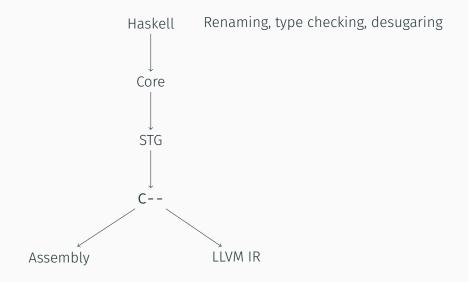
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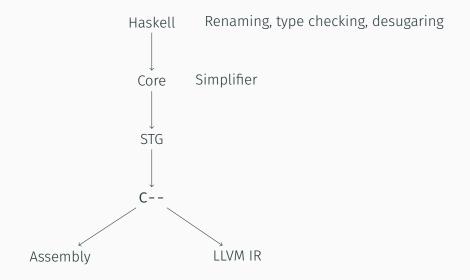
E.g. all primitive types but also **Array**# (which is not primitive)

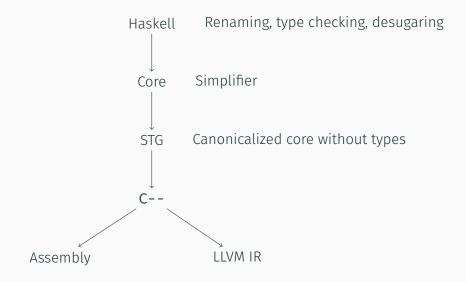
Definition of the **Int** type











Core's Expr Type

```
data Expr b
 = Var Id
   Lit Literal
   App (Expr b) (Arg b)
   Lam b (Expr b)
   Let (Bind b) (Expr b)
  | Case (Expr b) b Type [Alt b]
  | Cast (Expr b) Coercion
   Tick (Tickish Id) (Expr b)
  Type Type
  | Coercion Coercion
 deriving Data
```

Viewing Core

- · -ddump-simpl or -ddump-prep
- · Suppress info that you don't care about
 - · -dsuppress-idinfo
 - · -dsuppress-ticks
 - · -dsuppress-module-prefixes
 - · -dsuppress-all
- GHC plugin that outputs core as HTML

Mental Model for Core

let

Allocates a thunk on the heap

case

Forces evaluation to WHNF

Naive Sum

```
sum :: [Int] -> Int
sum [] = 0
sum (x : xs) = x + sum xs
```

Tail-Recursive Sum

```
sum :: [Int] -> Int
sum = go 0
   where
     go acc [] = acc
     go acc (x : xs) = go (x + acc) xs
```

Force the Accumulator

```
sum :: [Int] -> Int
sum = go 0
  where
    go acc [] = acc
    go acc (x : xs) =
       let acc' = x + acc
    in acc' `seq` go acc' xs
```

BangPatterns

```
{-# LANGUAGE BangPatterns #-}
sum :: [Int] -> Int
sum = go 0
  where
    go acc [] = acc
    go acc (x : xs) =
       let !acc' = x + acc
    in go acc' xs
```

WHNF

- seq only evaluates to WHNF
- Be careful with tuples!(x,y) `seq` ... will not evaluate x and y
- Use the deepseq lib for evaluating to NF

Strictness Annotations in Data Types

```
data Point = Point !Int !Int
```

Whenever you evaluate **Point** to WHNF, you also evaluate the two fields to WHNF.

Often easier to use than seq/BangPatterns

Avoiding Space Leaks

Rule of Thumb

Constant-size (e.g. Int) accumulators are often problematic

Detecting Space Leaks

- Limit the stack size +RTS -K\${n}K
- Get a stacktrace with +RTS -xc -K\${n}K

Specialization and Inlining

Specialization

- Specialize type parameters
- · Remove type class dictionaries

Inlining

· Inline definition at call site

Cross-Module Specialization and Inlining

- Specialization/Inlining only possible if definition (=unfolding) is available
- · Unfoldings of small definitions are automatically exposed
- {-# INLINABLE f #-} forces GHC to expose f's unfolding
- You might also want to expose unfoldings of definitions used by f

Specialization

- GHC will automatically try to specialize definitions at use-sites
- Create specializations using
 {-# SPECIALIZE f :: Int -> Int #-}
 - · Also creates specializations of functions called by f

- {-# INLINE f #-} makes GHC very eager to inline f
- · Use cautiously!
 - · Can blow up compile times significantly
 - · Can increase code size without bringing benefits

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- Use cautiously!
 - · Can blow up compile times significantly
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- Note: {-# INLINABLE f #-} does not make GHC more eager to inline f

The following two definitions are equivalent.

$$f a b = ...$$

$$f = \a b \dots$$

Or are they?

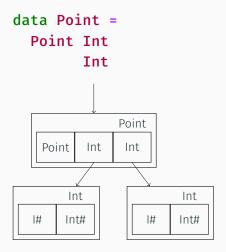
$$f a b = ...$$

$$f = \ad b \dots$$

Call Arity

GHC will only inline fully saturated function applications!

Controlling Memory Layout



Controlling Memory Layout

Automatic Unpacking

- GHC is quite good at automatic unpacking
- But only if it can detect that an argument is strict
- · Sometimes you need to help it

```
f :: Vector Int -> ...
f xs = ...
where n = Vector.length xs
```

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f :: Vector Int -> ...
f xs = ...
where !n = Vector.length xs
```

Continuation Passing Style

```
main :: IO ()
main =
  case loop2 100 (10, 10) of
    (x, y) \rightarrow print (x - y)
loop2 :: Int -> (Int,Int) -> (Int,Int)
loop2 n (x, v)
  | n > 0 = loop2 (n - 1) (x + 1, y - 1)
  | otherwise = (x, y)
```

Continuation Passing Style

Convert

into

$$f :: a \rightarrow (b \rightarrow r) \rightarrow r$$

Can avoid allocations and unnecessary case distinctions

Continuation Passing Style

Writing Your Own Optimizations: Rewrite Rules

map Fusion

$$map f . map g = map (f . g)$$

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map f . map g = map (f . g)
```

build/foldr Fusion

```
build
    :: (forall b. (a -> b -> b) -> b -> b)
    -> [a]
foldr :: (a -> b -> b) -> b -> [a] -> b

foldr f a (build g) = g f a
```

Writing your Own Optimizations: Rewrite Rules

- GHC does not check correctness of rules
- GHC does not check termination of rules
- · Use phases to control interaction of rules and inlining

Array Primitives in GHC

Array#

- Array of boxed values
- Card table to avoid having to scan unmodified entries in GC

SmallArray#

- Array of boxed values
- · No card table

ByteArray#

- Region of raw memory
- · Pinned and unpinned

High-Level Array Libraries

- primitive provides PrimArray wrapper around ByteArray#
- vector provides boxed, unboxed and Storable vectors
 - Fusion
 - Slicing

Basic Data Structures

- containers is mostly pretty good!
 - Use the specialized data structures for Int: IntSet and IntMap
- · unordered-containers has a fast, persistent HashMap
- Mutable hashtables from the hashtables package are often slower

Conclusion

- · GHC is impressively good at optimizing high-level code
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- GHC gives us the tools to control specific aspects of our programs
- · If all else fails, GHC has a great C FFI

More Information

- The Spineless Tagless G-machine
- Detecting Space Leaks
- Inlining and Specialisation
- · GHC User's Guide
- The GHC Commentary